Problem detection in real-time systems by trace analysis

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POLYTECHNIQUE MONTRÉAL

TRACE COMPASS

Outline

- Introduction
- Literature review
- Modeling
- Views
- Results
- Conclusion

Introduction : problematic

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Music player trace in Trace Compass

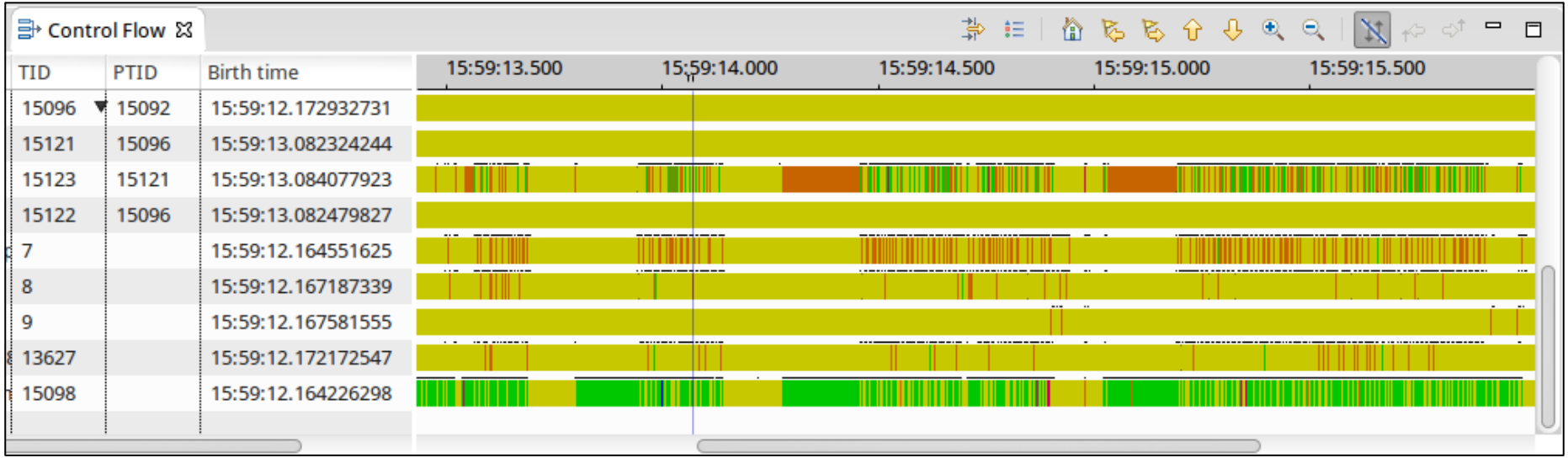

Figure 1 : Multiple executions of an audio player

Introduction : problematic

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Advantages of tracing real-time systems

- Low **overhead**
- Low **jitter**
- Access to **specific** information (priority, scheduling policy, etc.)

What is missing?

- Real-time **specific** user tools
- Show **useful** data

Introduction : goals

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1. Develop a **model** to define real-time task **executions** in a trace
2. Identify common **problems** in real-time systems and useful **information** to analyze them
3. Develop a method to analyze the **trace segment** corresponding to an execution to identify if the execution presents a **problem**

Introduction : definition

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- Real-time task : execution time, deadline, period (optional)
- Execution : periodic, sporadic
- Hard/soft real-time

PREEMPT_RT

- Priority inheritance for mutex in kernel
- Reduce non-preemptive sections in kernel

Scheduling policies

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- Normal
  - SCHED_OTHER : standard
  - SCHED_BATCH
  - SCHED_IDLE
- Real-time
  - SCHED_FIFO
  - SCHED_RR : with time quantum
  - SCHED_DEADLINE : Global Earliest Deadline First, highest user controllable priority

Scheduling policies

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- SCHED_FIFO and SCHED_RR
  - A deadline can be missed even if there was a valid scheduling to respect all deadlines
- SCHED_DEADLINE
  - No deadline will be missed if there is a valid scheduling

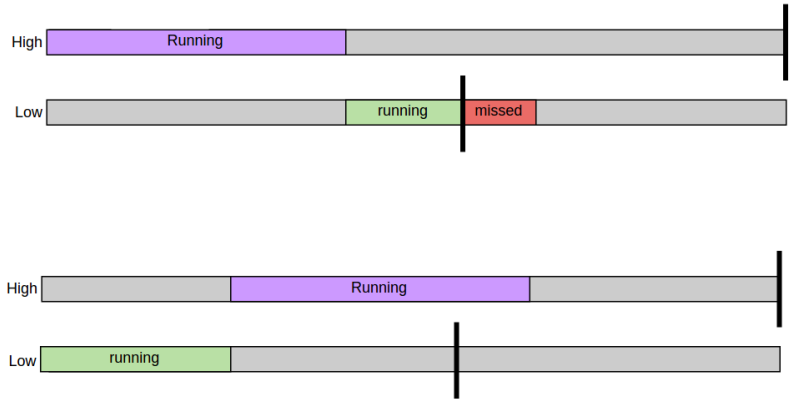

Figure 2 : Deadline missed

Scheduling policies

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- SCHED_FIFO and SCHED_RR
  - The highest priority task will always execute if it is able to
- SCHED_DEADLINE
  - If there is a missed deadline, it can be on a highest priority task (for the user, because there is no priority set)

Figure 3 : Highest priority

Priority inversion

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The high priority task is blocked by the low priority task that is preempted because the medium priority task is running.

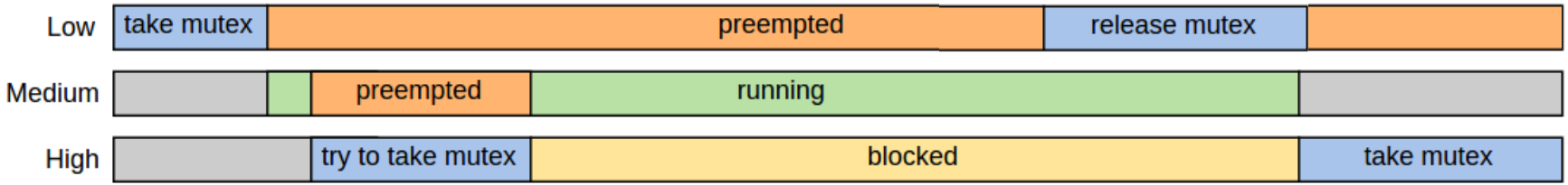

Figure 4 : Priority inversion

Priority inversion

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Priority ceiling protocol

- Better if the high priority task accesses the resource more often than the low priority task, because it is faster and has fewer context switches, but it can give an unnecessary high priority to the lower task

Figure 5 : Priority ceiling protocol

Priority inversion

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Priority inheritance

- Better if the low priority task accesses the resource more often

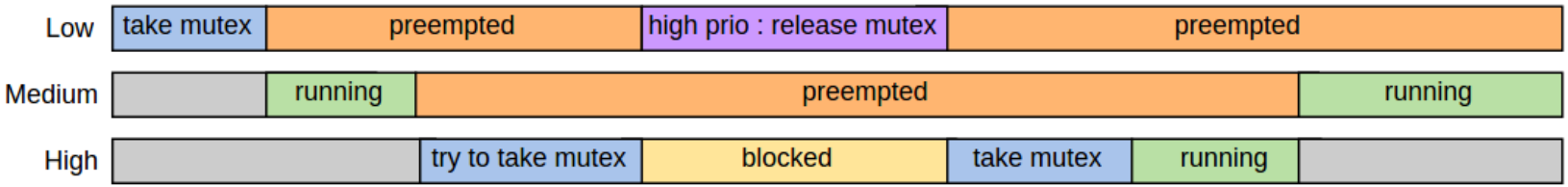

Figure 6 : Priority inheritance

Literature review

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Linux low-latency tracing for multicore hard real-time systems (Beamonte, 2013)

- LTTng-UST modification to **reduce** the added **latency**
- Demonstrated **low latency** tracing with LTTng

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Real-time Linux analysis using low-impact tracer (Rajotte, 2014)

- Recreate the task states using kernel events
- Compare executions of a task
- Sort the executions by running time
- Limitations
  - Threads need to have different priorities
  - Model is fixed
  - Not working with SCHED_DEADLINE
  - Manual analysis to find problems
  - Problems when more than one processor

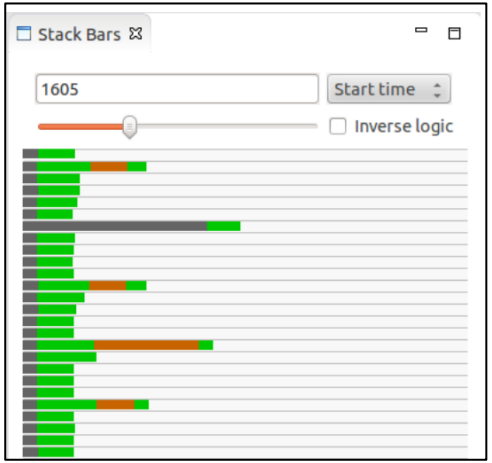

Figure 7 : Original stackbars view

Modeling

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Advantage of using only kernel events

- No need to modify the application source code to add tracepoints manually

Modeling

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- Identify executions automatically and then let the users choose between some valid models
  - Define a support ratio
  - Find all event types that are more frequent than the ratio
  - Increase the episode sizes using the fact that the sub-episodes must also be supported
  - Difficulties :
    - Using only event types
    - Execution time and memory usage
    - Many possible resulting models

Modeling : method

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State machine

- User identifies :
  - an execution or
  - events that define the start and the end (name, parameters with operations, etc.)
  - TIDs for start and end
  - Presets for common cases

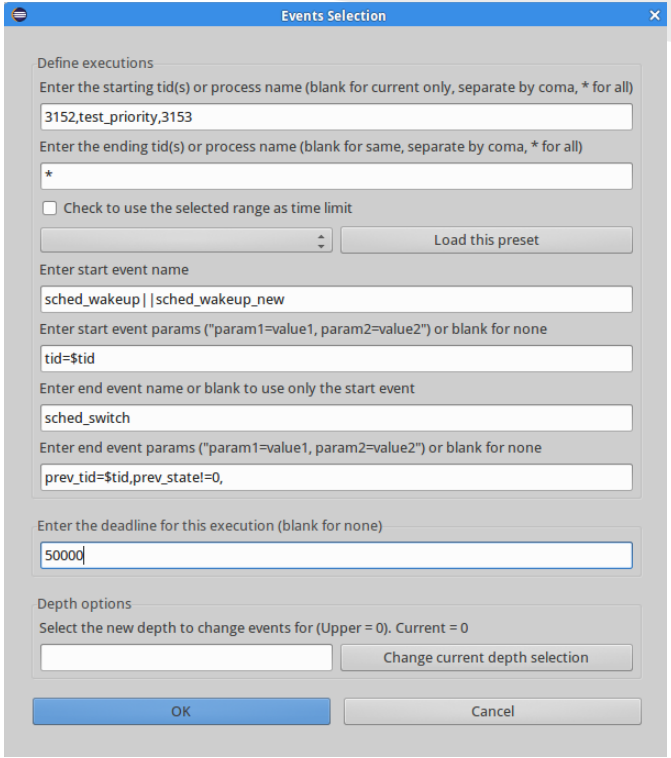

Figure 8 : Dialog to define model

Modeling : method

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State machine

- Remove execution
- Add execution
- Define an execution as invalid and recalculate
  - Will suggest some modifications to the model based on differences between valid and invalid executions
  - The user can select the ones he wants to apply

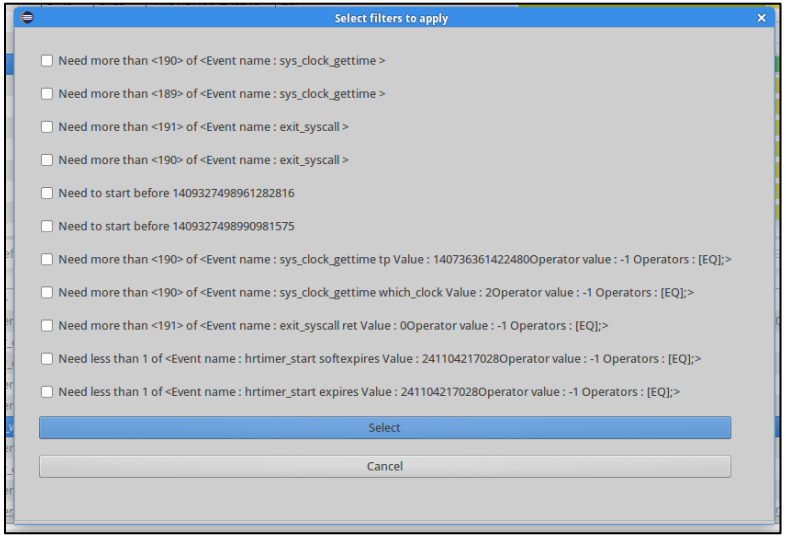
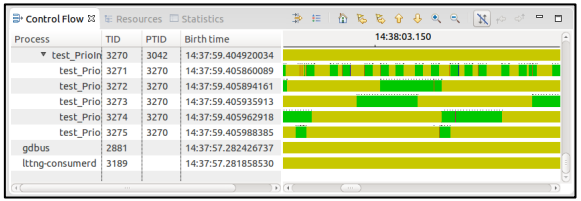


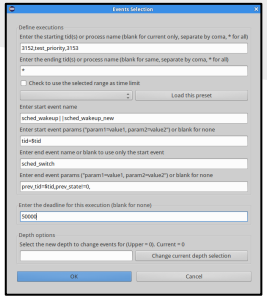
Figure 9 : Dialog to select modifications to apply

Overview

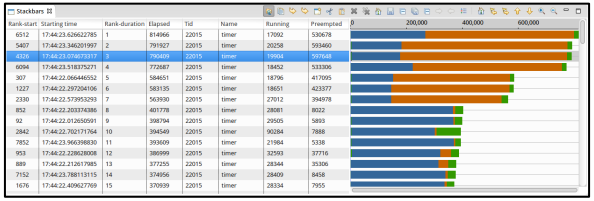
1) Control Flow View



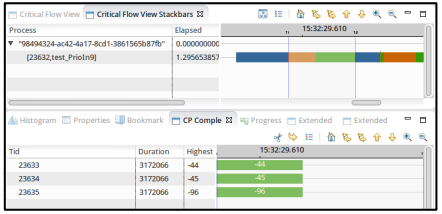
2) Define executions



3) Stackbars View



4) Critical Flow View with CP Complement view



5) Other views

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Stackbars view

Stackbars

Rank-start	Starting time	Rank-duration	Elapsed	Tid	Name	Running	Preempted
6512	17:44:23.626622785	1	814966	22015	timer	17092	530678
5407	17:44:23.346201997	2	791927	22015	timer	20258	593460
4326	17:44:23.074673317	3	790409	22015	timer	19904	597648
6094	17:44:23.518375271	4	772687	22015	timer	18452	533306
307	17:44:22.066446552	5	584651	22015	timer	18796	417095
1227	17:44:22.297204106	6	583135	22015	timer	18651	423377
2330	17:44:22.573953293	7	563930	22015	timer	27012	394978
852	17:44:22.203374386	8	401778	22015	timer	28081	8022
92	17:44:22.012650591	9	398794	22015	timer	29505	5893
2842	17:44:22.702171764	10	394549	22015	timer	90284	7888
7852	17:44:23.966398830	11	393609	22015	timer	21984	5338
953	17:44:22.228628008	12	386999	22015	timer	32593	37716
889	17:44:22.212617985	13	377255	22015	timer	28344	35306
7152	17:44:23.788113115	14	374956	22015	timer	28409	8458
1676	17:44:22.409627769	15	370939	22015	timer	28334	7955

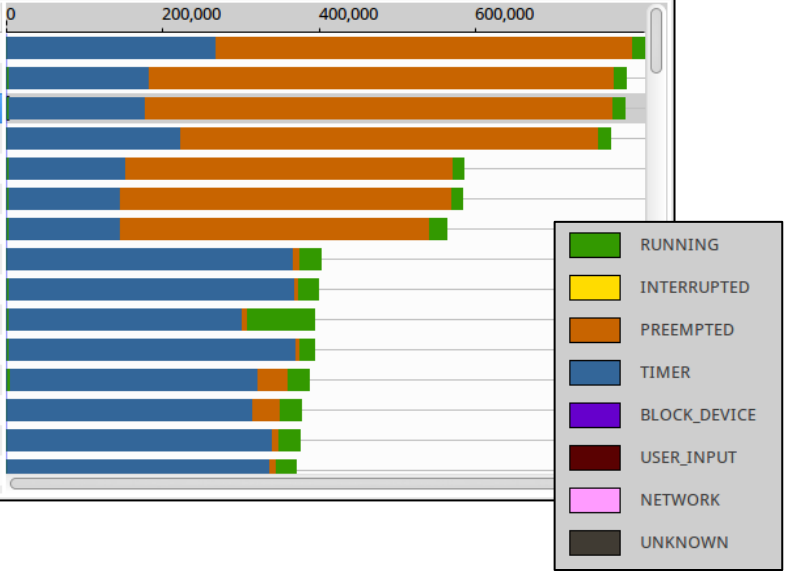

Figure 10 : Stackbars view

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- Supports
  - Thread pool
  - Nested executions

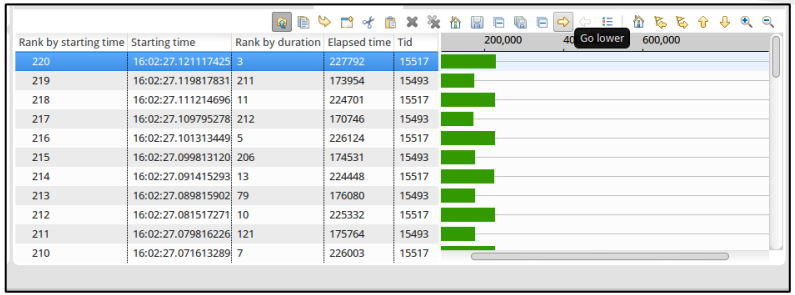

Figure 11 : Task on multiple threads

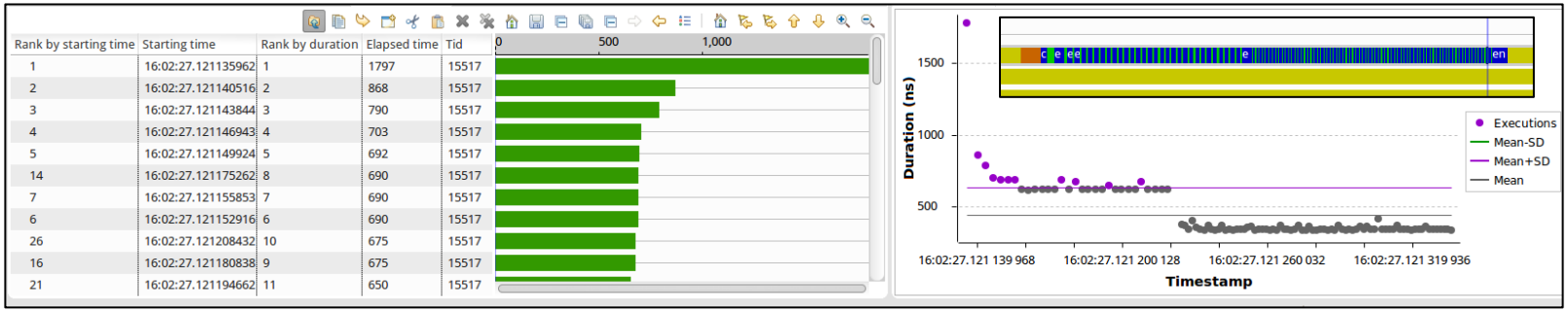

Figure 12 : Nested executions

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Time View

- View of duration by starting timestamp
- Synced with other views

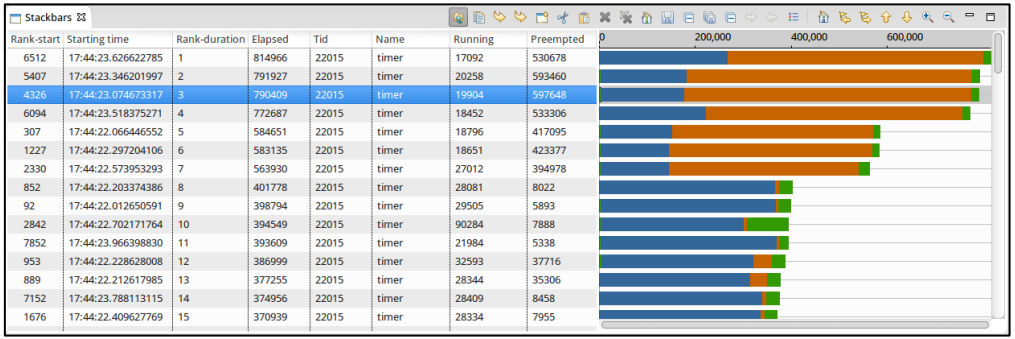

Figure 14 : Stackbars view

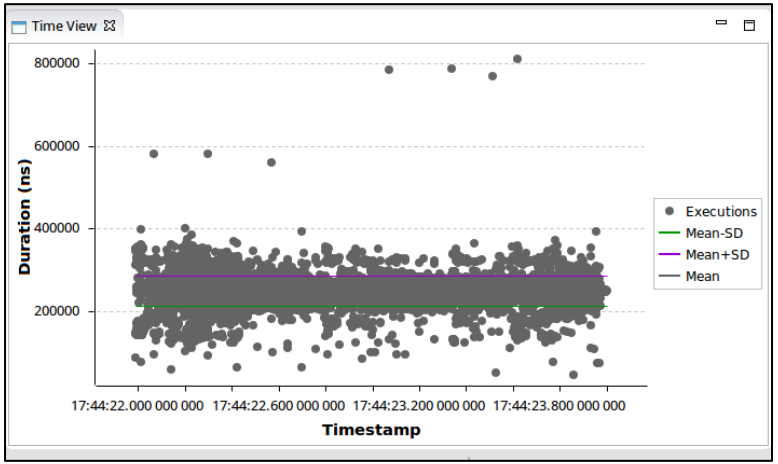

Figure 13 : Time view

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CP Complement View

- Show the priority of all running threads during preemption period of any thread in the critical path

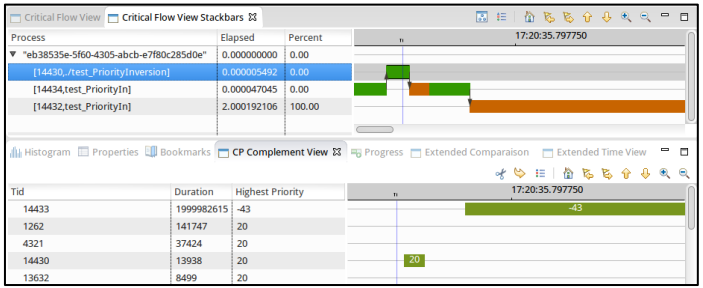

Figure 15 : CP Complement view

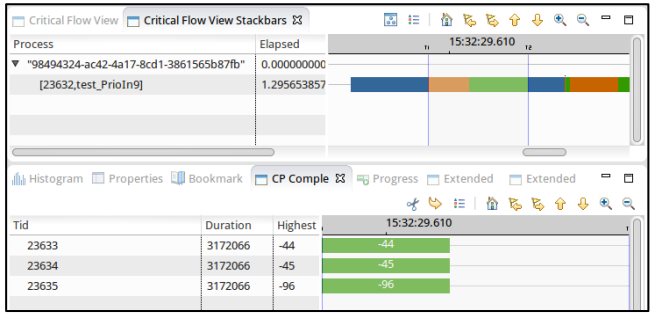

Figure 16 : CP Complement view

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CP Complement View

- Detect priority inversion

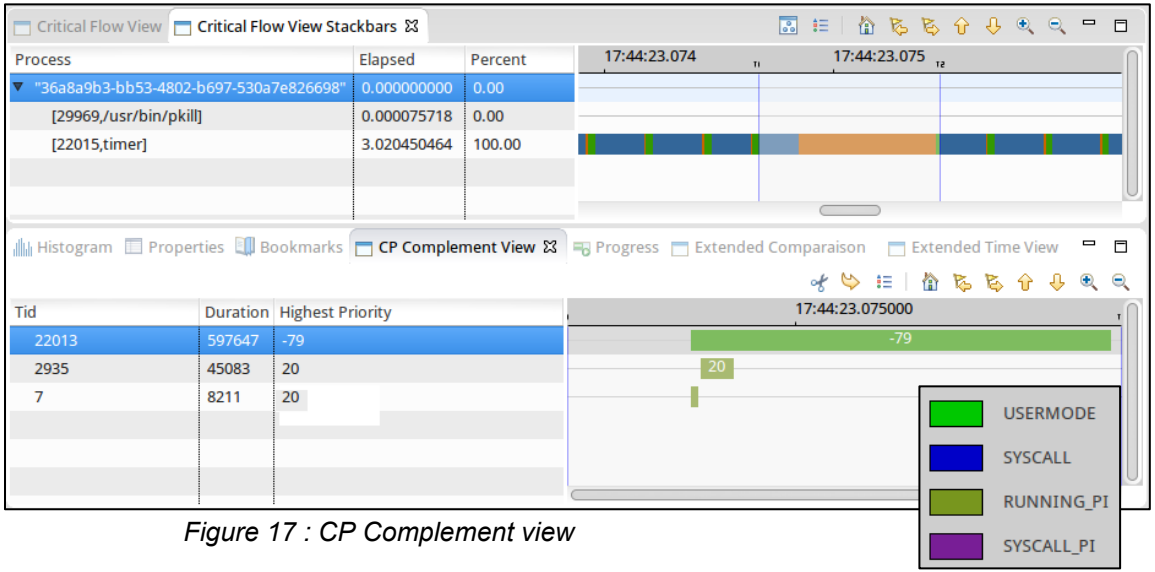
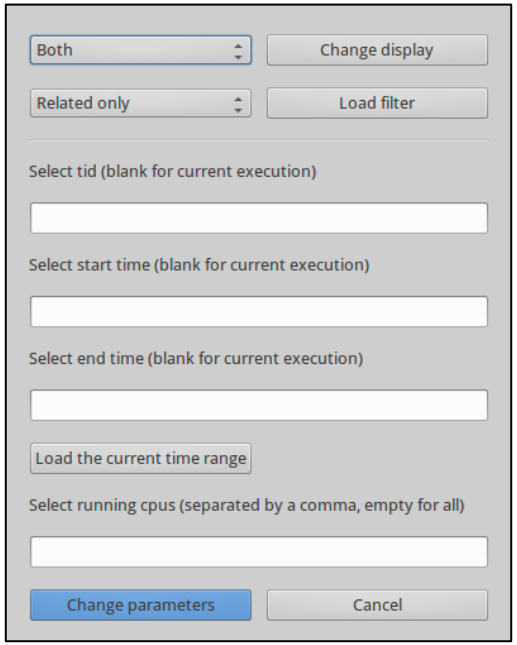


Figure 17 : CP Complement view



Example

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Find out why some executions take more time

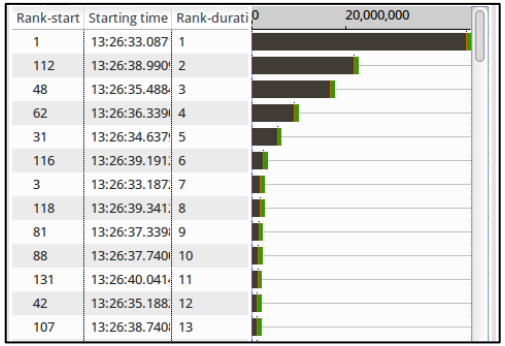

Figure 18 : Problematic executions

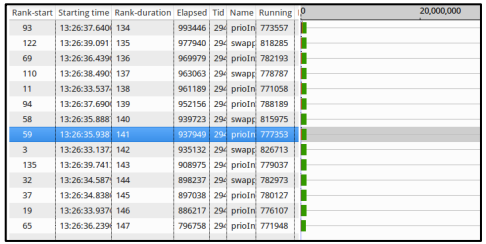

Figure 19 : Normal executions

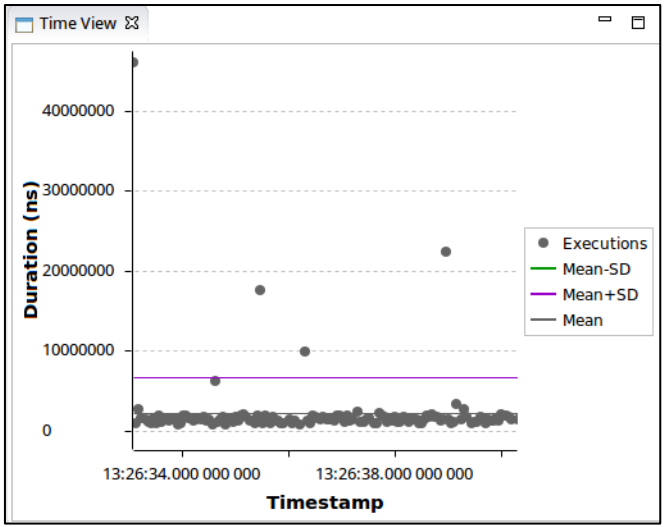

Figure 20 : Time View

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Normal execution

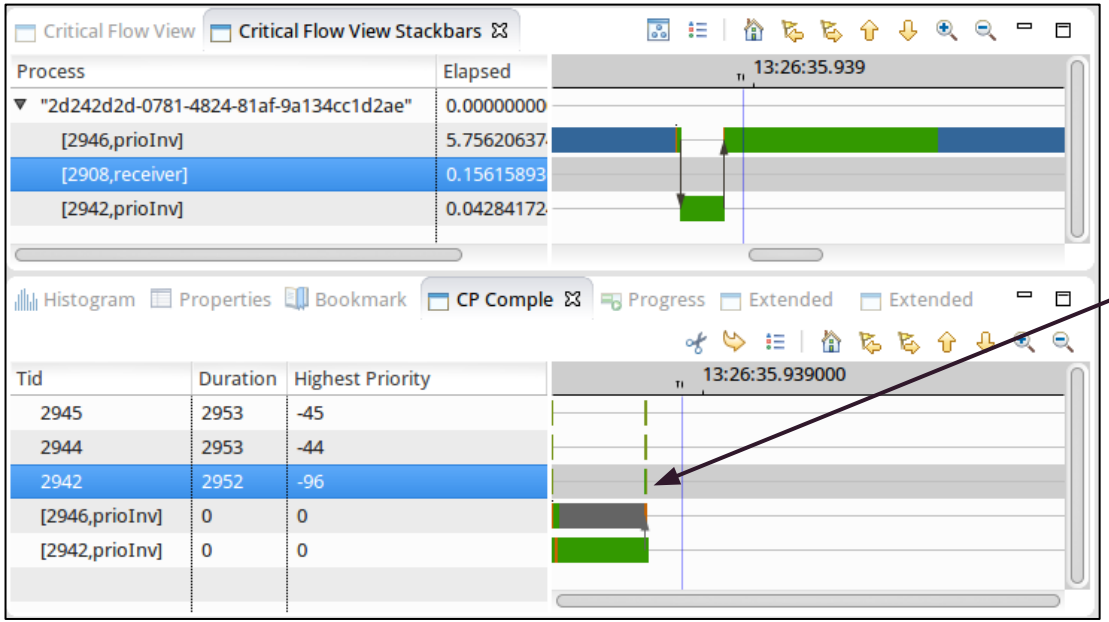

Figure 21 : CP Complement of a normal execution

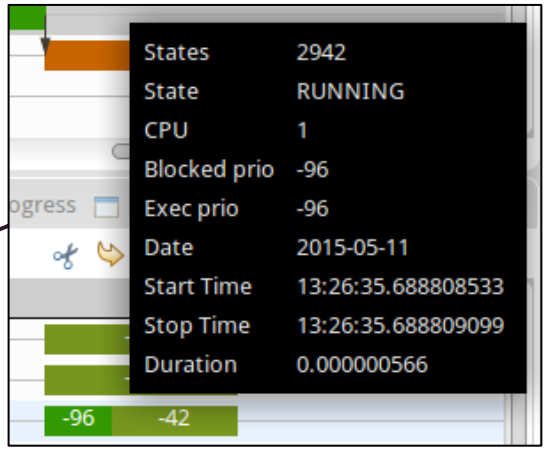

Figure 22 : There was priority inheritance

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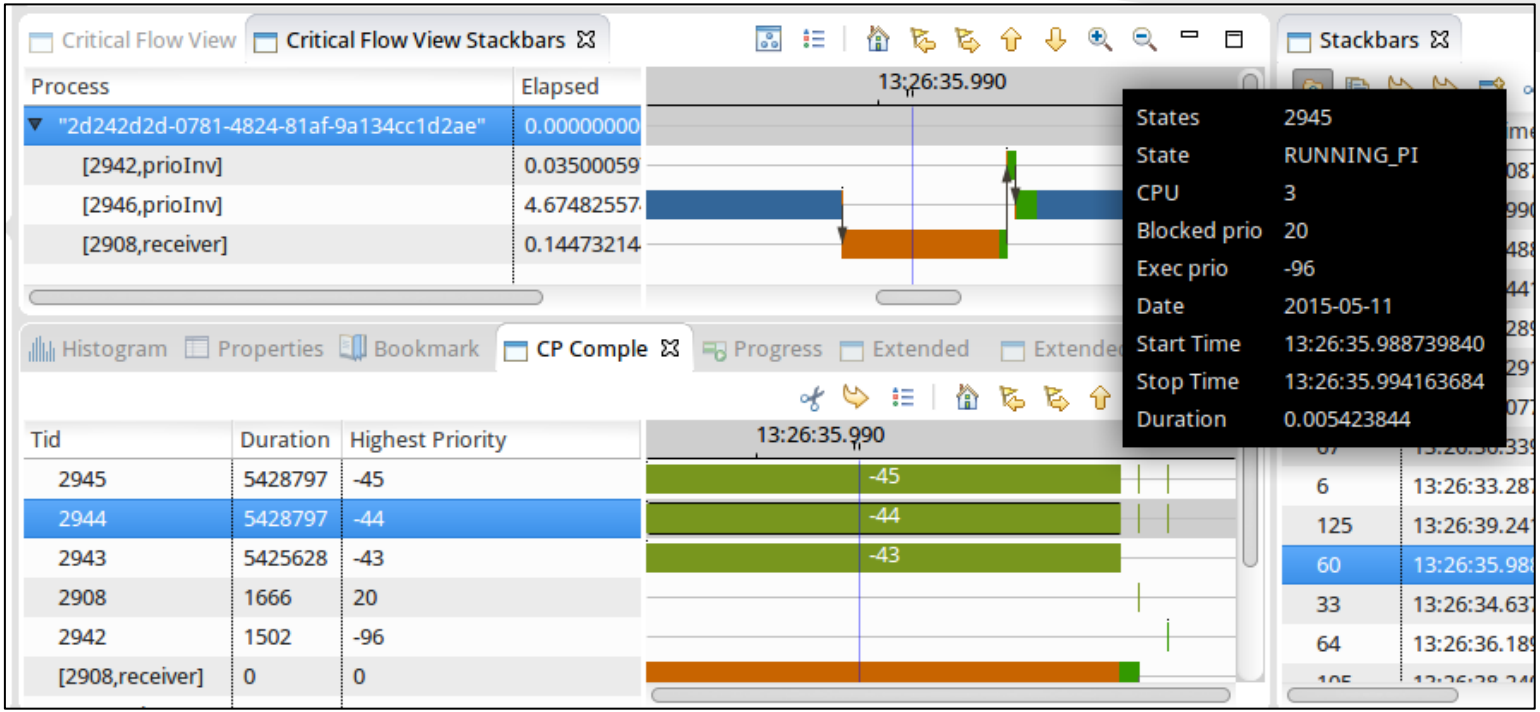

Figure 23 : CP Complement of a problematic execution

Other results

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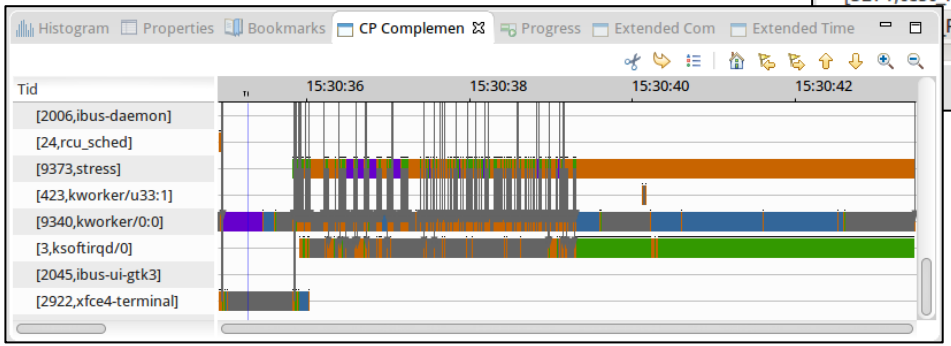
Output the dependencies during an execution

- priority
- directly related option



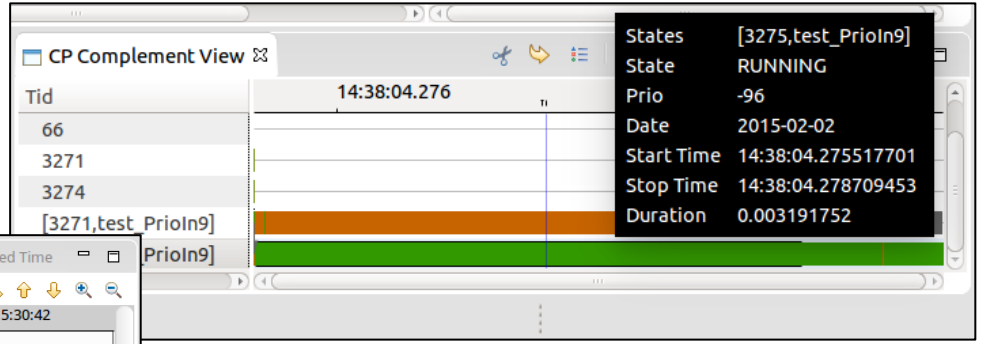


Figure 24 : CP Complement in related mode

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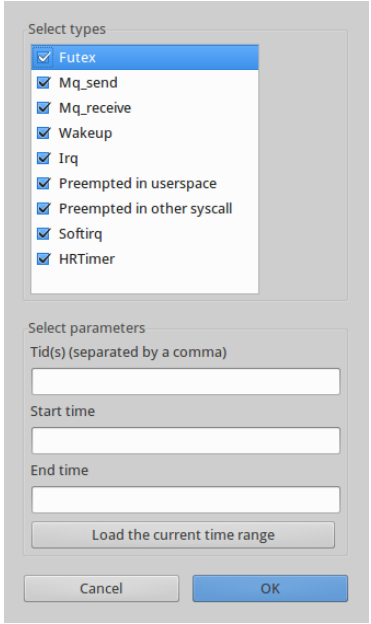
Extended comparison view



Type	Total running	Information
MQ_RECEIVE	5000077984	Queue=3
HRTIMER	5000054046	id = 18446612133080563168
MQ_RECEIVE	5000042860	Queue=3
HRTIMER	5000019151	id = 18446612133080563168

0 2,000,000,000 4,000,000,000

Histogram Properties Bookmark CP Comple Progress Extended Extended

Type	Total running	Information
OTHER	33130324	Event=clock_nanosleep
HRTIMER	33124087	id = 18446612135273365136
OTHER	32352676	Event=clock_nanosleep
HRTIMER	32345966	id = 18446612135273365136
OTHER	27422578	Event=clock_nanosleep
HRTIMER	27415884	id = 18446612135273365136
FUTEX	21833792	Futex=0x602140
FUTEX	16857471	Futex=0x602140
FUTEX	16079128	Futex=0x602140

0 20,000,000

Figure 25 : Extended comparison view

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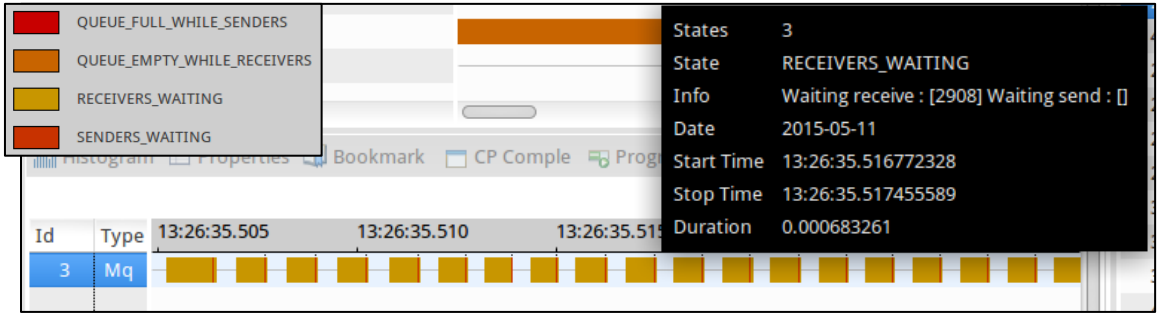
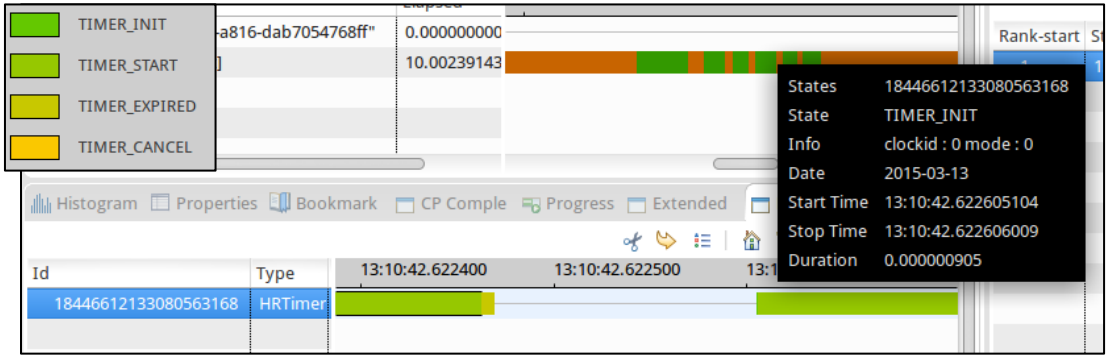
Extended time view : queue



HRTimer



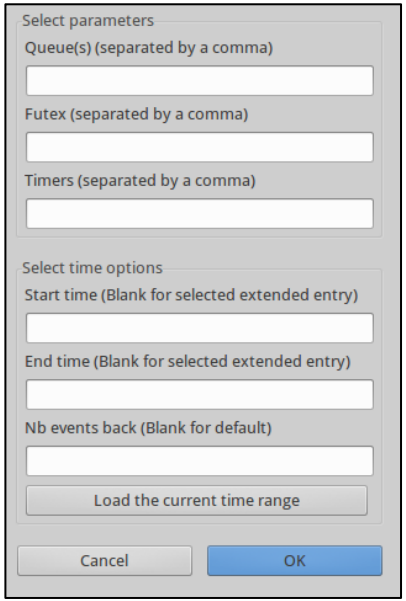


Figure 26 : Extended time view

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- Deadline analysis
  - Tell which executions missed their deadlines
  - User input

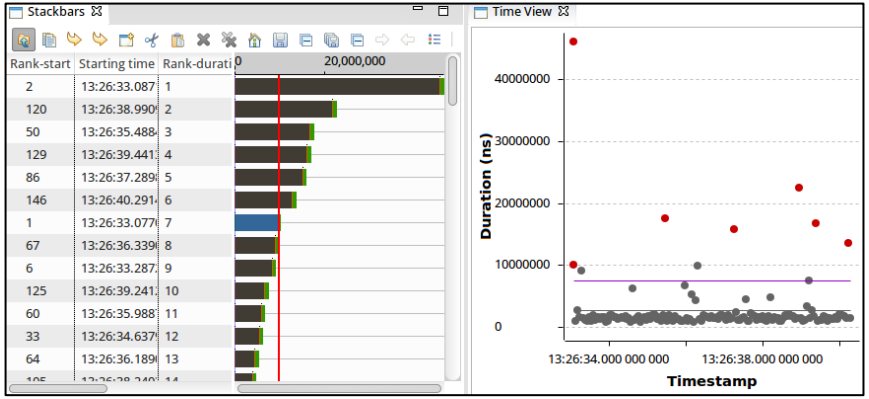

Figure 27 : Deadline

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- Future work
  - Modeling
    - Instrument complex real-time application in user-space and for each task, validate if it is possible to model only with kernel events
  - Analysis
    - Validate with real bugs
    - Add new analysis
- Questions?